

Detecting malicious documentation

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My research areas

- Practical security problems of Virtual Machine (VM)
 - Protect VM
 - Live memory forensic for VM
 - Malware scanner for VM
 - Leverage VM for various security-related areas
 - Malware analysis
 - Dynamic binary analysis
 - Vulnerability research
 - etc ...



Talk preview

- A new scanner named **D-Analyzer**, specially built to **detect malicious documentation files**
 - Reliably detect malicious attack from unknown documentation files
 - Independence of file formats
 - Zero-day attack detection is supported



Agenda

- Modern threats from malicious documentation
- **D-Analyzer** solution
 - Approach
 - Background on Tainting analysis
 - Architecture, Design & Implementation
- Discussions
- Conclusions
- Q & A



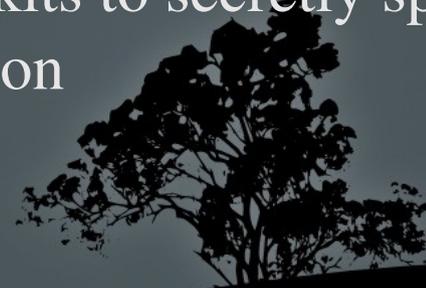
Traditional malware

- Executable file format
 - Mostly **.EXE** files
 - Trick users to run malware
 - Drive-by download
 - Social engineering
- Detect and clean with traditional malware scanner
 - Scan on-demand
 - Scan on-the-fly



Upcoming malware

- From **non-executable** documentation files
 - Attack embedded in any type of documentation files
 - .PDF, .DOC, .XLS, .PPT, .AVI, ...
 - Silently exploit vulnerable application when opening
 - Function under the same account privilege with current user
 - Privilege escalation is possible with other local vulnerabilities
 - Setup backdoors
 - Download & install trojan, rootkits to secretly spy on the victim, and steal information



Upcoming malware (2)

- How is it possible?
 - Most file formats are complicated, with lots of optional data fields
 - Handling file data properly is far from trivial, thus the implementation is vulnerable to security bugs
 - Buffer overflow, Format string, Integer overflow
- How to exploit?
 - Make victim to open a craft documentation file to trigger vulnerable code in vulnerable application
 - Embed malicious payload executed to exploit the vulnerability



Upcoming malware (3)

- Documentation-based attack become extremely popular recently
 - Not in theory, but getting more and more serious
 - Dominate other kind of bugs, and frequently published with PoC code
 - Why?
 - Client-side attack become a major trend
 - Finding client-side vulnerabilities is easier than server-based ones
 - Very few users are aware of this threat!
 - More users → more victims → generate more money for cybercrime
 - Make targeted attack feasible
- 

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MOAUB – 30 days of 0days, Binary Analysis and PoCs



17th August 2010 - by admin

The **Abysssec Security Team** is about to unleash its **Month Of Abysssec Undisclosed Bugs** on us. Starting on the 1st of September, Abysssec will release a collection of 0days, web application vulnerabilities, and detailed binary analysis (and pocs) for recently released advisories by vendors such as **Microsoft, Mozilla, Sun, Apple, Adobe, HP, Novel**, etc. The 0day collection includes PoCs and Exploits for **Microsoft Excel, Internet Explorer, Microsoft codecs, Cpanel** and others. The MOAUB will be hosted on the Exploit Database, and will be updated on a daily basis. Get your hard-hats on, your VM's and debugging tools organized – it's gonna be a an intensive ride. Follow both the **exploit-db** and **Abysssec** twitter feed to keep updated!



Sample vulnerabilities (1)

- **CVE 2010-3000**: Multiple integer overflows in the ParseKnownType function in RealNetworks RealPlayer 11.0 through 11.1 and RealPlayer SP 1.0 through 1.1.4 on Windows allow remote attackers to **execute arbitrary code** in an **.FLV** file.
- **CVE-2010-1900**: Microsoft Office Word 2002 SP3, 2003 SP3, and 2007 SP2; Microsoft Office 2004 and 2008 for Mac; Open XML File Format Converter for Mac; Office Word Viewer; Office Compatibility Pack for Word, Excel, and PowerPoint 2007 File Formats SP2; and Works 9 do not properly handle malformed records in a **Word file (.DOC)**, which allows remote attackers to **execute arbitrary code** or cause a denial of service (memory corruption) via a crafted file, aka "Word Record Parsing Vulnerability."



Sample vulnerabilities (2)

- **CVE-2010-1246**: Stack-based buffer overflow in Microsoft Office Excel 2002 SP3 allows remote attackers to **execute arbitrary code** via **an Excel file (.XLS)** with a malformed RTD (0x813) record, aka "Excel RTD Memory Corruption Vulnerability."
 - **CVE-2010-1681**: Buffer overflow in VISIODWG.DLL before 10.0.6880.4 in Microsoft Office Visio allows user-assisted remote attackers to **execute arbitrary code** via a crafted **.DXF** file, a different vulnerability than CVE-2010-0254 and CVE-2010-0256.
 - **CVE-2008-2238**: Multiple integer overflows in OpenOffice.org (OOo) 2.x before 2.4.2 allow remote attackers to **execute arbitrary code** via crafted EMR records in an **.EMF** file associated with a StarOffice/StarSuite document, which trigger a heap-based buffer overflow.
- 

Sample vulnerabilities (3)

- **CVE-2010-0188**: Unspecified vulnerability in Adobe Reader and Acrobat 8.x before 8.2.1 and 9.x before 9.3.1 allows attackers to cause a denial of service (application crash) or possibly **execute arbitrary code** via unknown vectors (**.PDF** file).
- Media Player Classic v6.4.9.1: **Buffer Overflow Exploit** for **.AVI** file
<http://www.exploit-db.com/exploits/11535/>
- Corel VideoStudio Pro is prone to a **remote buffer-overflow vulnerability** because the software fails to perform adequate boundary checks on user-supplied data (**.MP4** file)
<http://www.securityfocus.com/bid/40963/discuss>
- **CVE-2010-0718**: Buffer overflow in Microsoft Windows Media Player 9 and 11.0.5721.5145 allows **remote attackers to attack** via a crafted **.MPG** file

Sample vulnerabilities (4)

- **CVE-2008-4434**: Stack-based buffer overflow in (1) uTorrent 1.7.7 build 8179 and earlier and (2) BitTorrent 6.0.3 build 8642 and earlier allows remote attackers to cause a denial of service (crash) and possibly **execute arbitrary code** via a long Created By field in a **.TORRENT** file.
- Media Player Classic - v 1.3.1774.0 (**.RM** file) Buffer Overflow
<http://www.exploit-db.com/exploits/12704/>
- **CVE-2009-4195**: Buffer overflow in Adobe Illustrator CS4 14.0.0, CS3 13.0.3 and earlier, and CS3 13.0.0 allows remote attackers to **execute arbitrary code** via a long DSC comment in an Encapsulated PostScript (**.EPS**) file.
- **OSVDB-ID 64984**: Easyzip 2000 v3.5 (.zip) 0day **stack buffer overflow PoC exploit**. When the application receives a malicious '**.ZIP**' file it fails to properly sanitize the 'filename' section on the zip resulting in a stack based buffer overflow.

Demo



Detect malicious documentations

- Not-yet the target of current malware scanners :-(
 - Mostly focus on **EXE** files
 - Or can only understand and detect malicious files in popular formats (**PDF, DOC, ...**)
 - But any file format can be vulnerable :-(
 - .FLV, .XLS, .PPT, .DXF, .EMF, .XML, .PDF, .WAV, .OGG, .MP4, .MPG, .RM, .ZIP, .TORRENT, .EPS, ...
 - Zero-day exploitation?



A dream scanner + analyzer

- Can detect malicious attack from documentation files
 - Support all kind of file formats
- Can report if a particular version of application is vulnerable or not
- Can analyze the attack behaviour
 - How the exploitation work?
 - What the exploitation does?
 - Modify system to do something evil?
 - Download + install malware?
 - Setup backdoors?
 - Steal information?



Assumption

- Given a **random** file of a **random format**, how can we know:
 - Is it a malicious file?
 - Which application and versions are vulnerable?
 - How the exploitation work?
 - What the exploitation does?
 - Modify system to do something evil?
 - Download + install malware?
 - Setup backdoors?
 - Steal information?



Idea?



Approach

- Open a suspected file with corresponding application
 - For ex, .PDF file can be open with Adobe Reader (with any interested version) or Foxit (any interested version)
- Detect if the suspected file exploit its application
- Monitor the exploitation behavior
 - To understand its exploitation process
 - To monitor the exploitation behavior to understand what it does at post-exploit stage



D-Analyzer solution

- Run everything inside a VM
 - Open a given file with related application
- Perform **dynamic tainting analysis** on VM to detect exploitation
- **Instrument** the VM to **monitor** vulnerable application and understand exploitation behavior



Dynamic tainting analysis

- "In order for an attacker to change the execution of a program illegitimately, he must cause a value that is normally derived from a trusted source to instead be **derived from his own input**" (James Newsome et al, Usenix Security 2005)
- **Tainted data**: data come from untrusted input
- Tracking flow of tainted data offers multiple advantages
 - Understand in detail **how data is (ilegally) used**
 - **Detect unknown (Zero-day) attack is possible**

Basic concepts

- **Tainted source**: data originated from untrusted input
- **Taint propagation**: monitor data flow, and mark related data as tainted
- **Taint sink**: tainted data is consumed in illegal way?
 - detect when data is illegally used, not written



Tainting propagation

- Two taint targets
 - Memory
 - Area of memory is tainted or not
 - Represented by start memory addr & range size
 - Hardware register
 - Regular registers: EAX, EBX, ECX, EDX, ESI, EDI, ESP, EBP
 - Other registers (segment registers, control registers, ...) can be ignored



Sample tainting propagation

- **Moving data insn** (MOV, XCHG, ...)
 - MOV EAX, EBX
 - EBX is tainted → EAX becomes tainted
 - EBX is untainted → EAX becomes untainted
- **Algorithm insn** (ADD, SUB, XOR, ...)
 - ADD EAX, EBX
 - EAX becomes tainted if EBX is tainted



Sample tainting propagation (2)

- "Constant" insn
 - `XOR EAX, EAX`
 - `SUB EAX, EAX`
 - Always untaint EAX, no matter what
- Algorithm insn (`ADD`, `SUB`, `XOR`, ...)
 - `ADD EAX, EBX`
 - EAX becomes tainted if EBX is tainted



Sample taint sinks

- `Jmp` & `Call` insn
 - `Jmp` or `Call` to tainted memory/register?
 - `EIP` is tainted!
- Format string functions
 - Format string argument is tainted, with dangerous format specifier like `%n` ?
- Sensitive system functions (like `WinExec`, `CreateProcessA`, `CreateRemoteThread`, ...)
 - Function arguments are tainted?



Tainting policy

- Taint source?
 - Network data, File, memory ranges, keyboard, ... ?
 - Taint propagation?
 - In a particular case, taint corresponding memory/register or not?
 - Fine-grain level of taint tracking:
 - bit/byte/word/dword, page level?
 - Taint sinks?
 - Where to have taint sinks?
 - What to do when tainted data is used at taint sink?
- 

Implicit data flow problem

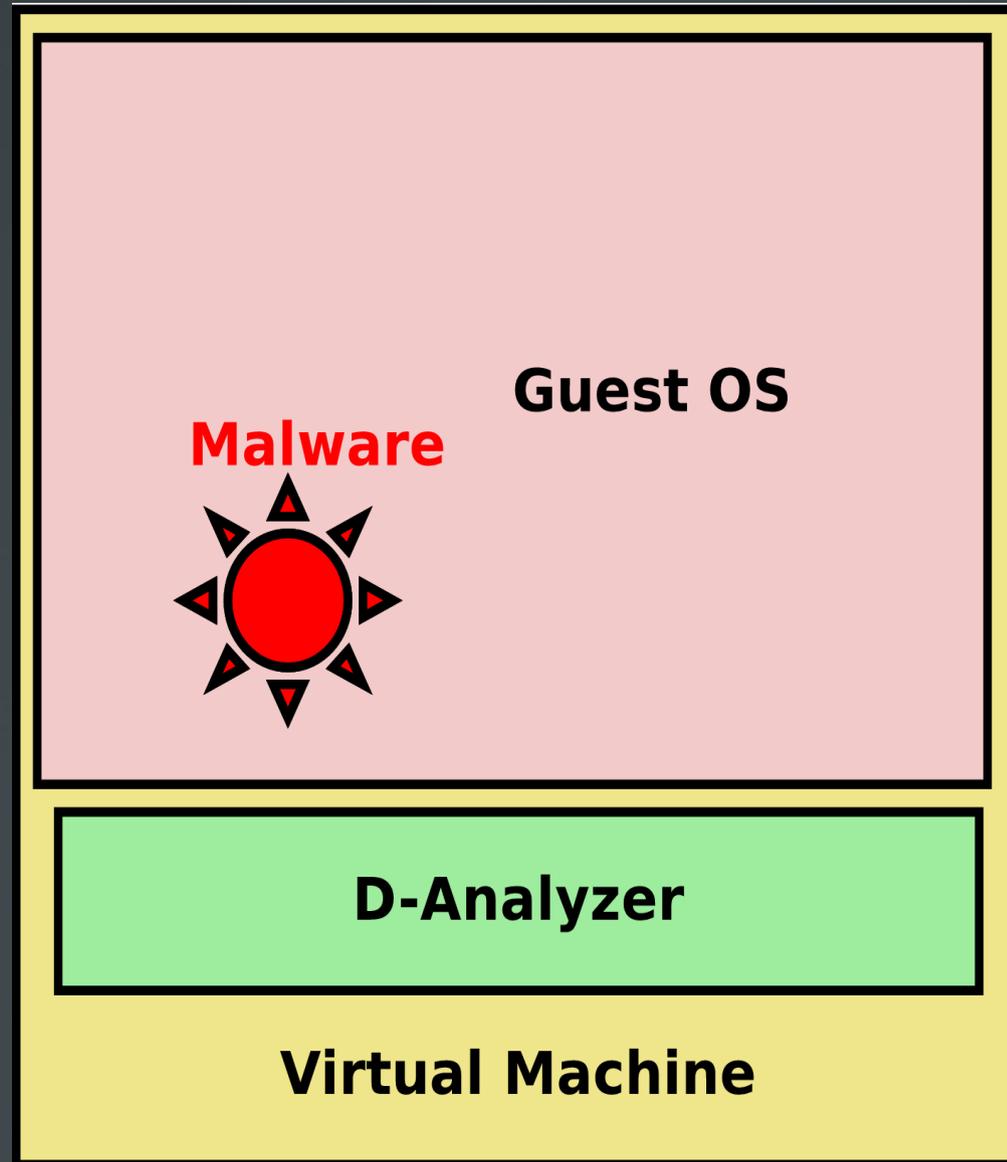
- Result data is **indirectly** influenced by tainted data
 - Data assignment based on branch insn

```
if (a == 1) b = 1;
else if (a == 2) b = 2;
```
 - Tainted data is used as index of an array (widely used in table lookup)

```
c = string[x];
```
- A main source of false negative issue → evade tainting analysis



D-Analyzer architecture



Solve the proposed questions

- **D-Analyzer** can detect exploitation thanks to tainting analysis mechanism
- **D-Analyzer** can report vulnerable applications, and also their versions
 - **Open suspected file with each application**
- Understanding the exploitation process and exploitation behavior is possible thanks to whole-system instrumentation monitoring



Add-on benefits

- **D-Analyzer** is (mostly) invisible to exploitation, so stealthily monitoring the exploitation process is feasible
 - Never run anything inside the VM
 - Perform instrumentation from below VM allows us to monitor everything, even ring-0 code
 - Analyzing kernel rootkit is possible



D-Analyzer requirement

- Understand VM context from outside
- Instrument guest VM execution
 - Intercept execution anywhere/anytime
- Access to VM context
 - Read/write to VM memory
 - Read/write CPU status
- Performing tainting analysis on suspected application



Understand VM context

- Must be done from outside, without any support of guest VM
 - VM introspection problem
 - Extract OS semantic objects from VM's memory
 - Support Windows OS

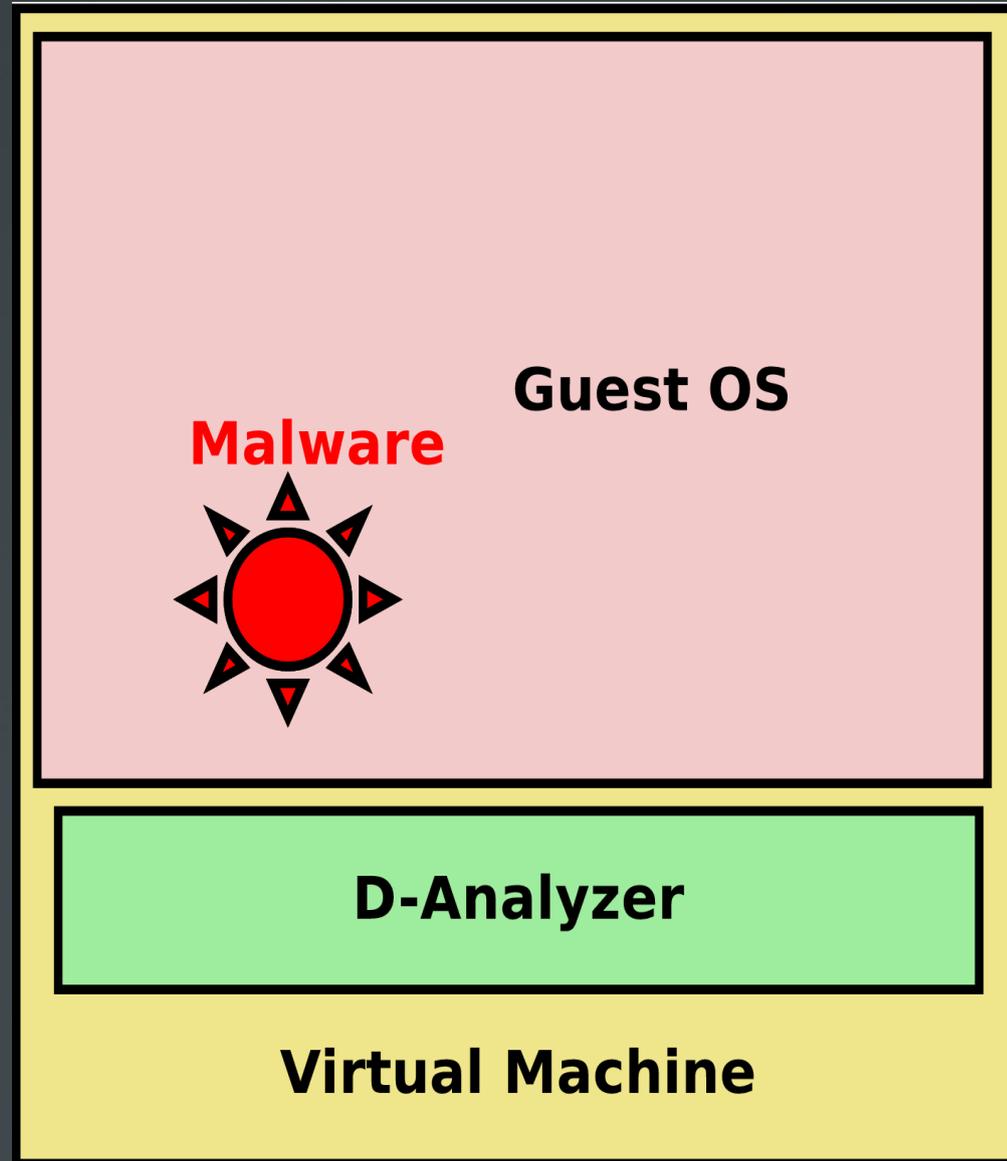


VM for D-Analyzer

- Choose VM for **D-Analyzer**
 - Open source, so customizable (therefore VMWare is not suitable)
 - Xen? KVM?
 - VirtualBox?
 - Bochs?
 - **Qemu?**
 - Version 0.12.5



D-Analyzer architecture



Instrument guest VM

- **Kobuta** framework
 - Generic instrumentation framework
 - Not only for **D-Analyzer**, but other future projects
 - Instrument QEMU's **dynamic binary translation** process
 - Put hooks at right places to call out to external handlers
 - Support dynamic loaded module built on top of **Kobuta**
 - Module provides instrumentation handlers to be executed when called from hooks



Instrument guest VM - Challenges

- Originally, QEMU provides **no support** for instrumentation
 - We are on our own, and have to build instrumentation framework from scratch
- QEMU uses **Just-in-time (JIT) compiler** to perform binary translation
 - Translated code is cached, and is not translated again if available in cache
 - We have to dig deeply into the translation process of QEMU to provide instrumental hooks



QEMU JIT compiler

- Translate guest code to **TCG IR**, then translate **TCG bytecode** to native (host) code to execute on host
- The translated code is **cached** to be reused (to improve performance)
- Translation is done on code block basis
- **CPU context** (register values) is saved at the end of each block
 - So **CPU context** is **synchronized** at **beginning** of each **block**
 - At **middle** of block, **CPU context** is **out-of-synch**



Kobuta framework

- Hooking various events useful for generic purposes
 - Fine grain instrumentation
 - Begin/end of instruction/block
 - Jump/call insn
 - Interrupt begin/end
 - Sysenter/Sysexit/Syscall/Sysret
 - Input/Output insn
 - Update control registers (CR0, CR3, CR4)
 - RDMSR, WRMSR (read/write to Model-Specific-Register)
 - Memory access (read/write)
- 

Kobuta module

- Need to register with **Kobuta** framework for interested instrumentation events
 - Then **provide instrumentation handlers** for those events
 - Handlers be executed when events happen in guest VM
 - Leverage **exported functions** (from **Kobuta** framework) to manage guest VM from module
 - Pause and Resume VM on demand
 - Read and write to VM's memory (physical & virtual memory and CPU context
 - Single step mode
 - Dynamically enable/disable instrumentation events
- 

Kobuta module

- Design Kobuta module to be just a Dynamic Linked module
 - .so file in Linux, .DLL file in Windows
 - Loadable into Qemu process, and supported by OS services
 - Easy to implement your Kobuta module (just a normal DL module running in host OS)



Manage Kobuta module

- **Load module** into Qemu process
 - Simply using DLL service provided by host OS
 - `dlopen()` in Linux
- **Unload module** from Qemu process
 - Also use DLL service of host OS
 - `dlclose()` in Linux
 - But how about code still running?



Unloading Kobuta module

- Use **reference counter** for Kobuta instrumentation handlers
 - **Increase** counter before running a handler, and **decrease** it when done
 - Only run a handler when its module is in enable state
 - Have a manage thread to unload Kobuta module
 - Put module in disable state
 - Periodically checking for ref counter, and unload module when $\text{refcount} = 0$



D-Analyzer tainting analysis

- Taint source
 - A **suspected file**, open with related application
- Taint sinks
 - **Jmp/Call** target
 - Detect Buffer Overflow exploitation
 - **Format string functions**
 - Detect Format String exploitation
 - Sensitive functions (such as **WinExec**, **CreateProcessA**, **CreateRemoteThread**, ...)
 - Detect false negatives

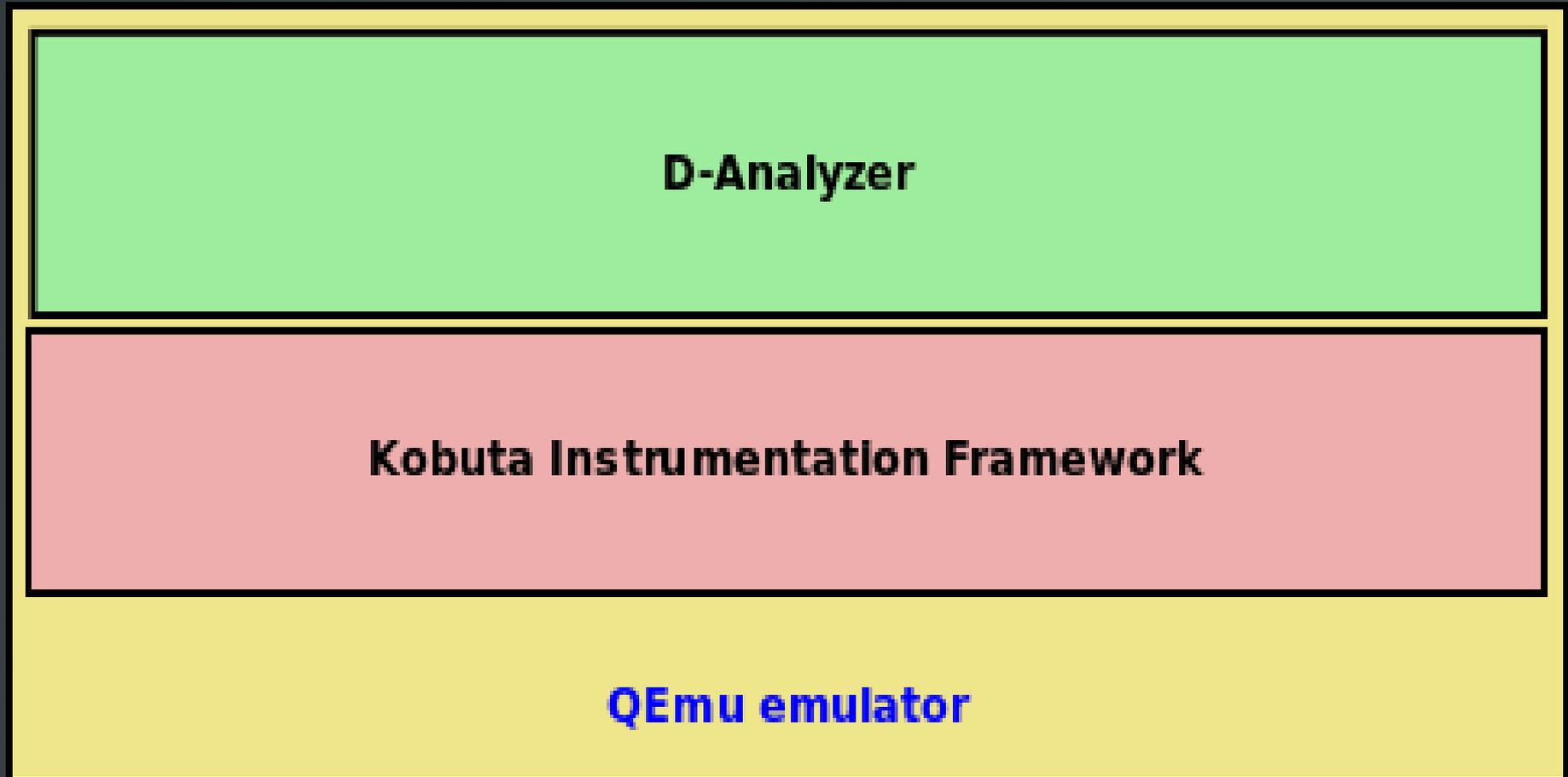


D-Analyzer tainting analysis (2)

- Taint propagation
 - Usual policy for moving data insn, algorithm insn
 - Tainting **table lookup result** if **index register** is tainted (implicit data flow tracking problem)
 - Tainting output data by **modeling function logics**



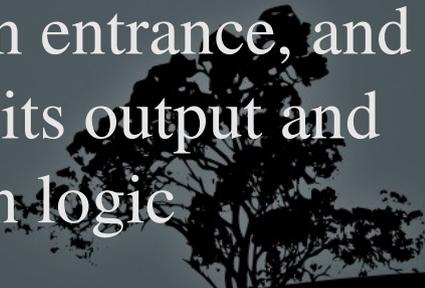
D-Analyzer module architecture



D-Analyzer implementation

- A **Kobuta** module
 - Register several instrumentation events
 - **JumpCall**
 - To capture interested function execution & their parameters
 - To handle sinks as Jump/Call targets
 - **Begin/end insn**
 - To perform tainting analysis on instruction level
 - **Update CR3**
 - To know if we are switching to our monitored processes
- 

D-Analyzer implementation (2)

- Employ several tactics to speed up tainting analysis process
 - On-demand instrumentation
 - Turn off insn level instrumentation when switching to other processes
 - Turn off insn level instrumentation on "uninteresting" functions, which produce no output nor modify parameters
 - On-demand tainting analysis by modeling function logic (such as string functions or File related functions)
 - Stop tracking functions at function entrance, and resume when function returns. Tainting its output and parameters according to function logic
- 

D-Analyzer implementation (3)

- Tracking data flow from **given input file**
 - Open (tainted) input file with the related application
- Report if **tainted data is illegally used at sinks**
 - Conclude **this file is malicious** → report
 - Also this particular application (of this particular version, too) is **vulnerable with this malicious file**
- Repeat the checking process with **all the requested applications**, on the **same input file**



Discussion

- **Implicit data flow** is not a major problem for us
 - Can be mostly handled by tainting public functions
 - Unlike malware, legitimate applications never try to evade tainting analysis :-)
 - False negative is greatly reduced
- **Control flow corruption problem** is solved, but exploitation targeting **data & data pointers** remain a problem
 - Partly solved by monitoring sensitive functions



Future plan - Development

- Improve performance
 - Other tricks to speed up taiting analysis
 - Using **KVM** to speed up emulation even further
 - Even currently, **KQEMU** is not too bad, either
- Track down the **culprit area** in malicious input file
- Focus on exploitation corrupting **data & data pointers**
- **Extract shellcode and analyze exploitation behavior**



Related works

- BitBlaze project
 - Based on old QEMU version (0.9.1)
 - Too slow due to whole system taint tracking
 - File input as tainted source does not seem to be supported
- Joedoc.org
 - Limited documentation files and applications
 - Signature-based detection
- Vichack.ca
 - Unknown technique



Conclusions

- **D-Analyzer** is an effective scanner for malicious documentations
 - Support all kind of file formats
 - Can detect zero-day malicious attacks
 - Analyze malicious attacks is possible

