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Automatic Malware Analysis

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Malware

Program which infiltrate machine with out the user awareness.

Few Vectors:

Exploits - CVE-2008-4250 (MS-RPC Vulnerability) Conficker.

Social Engineering – email , phishing.

Drive by Downloads – infected web site

0days – new vulnerabilities , RBN.

can install keyloggers to steal sensitive information

can hijack machines, install applications, servers, filter drivers

can make the machine participate in DDoS Attacks



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The Challenge ..

Analysis Facts

Malwares authors tends to use packers.

(Originally used to reduce the binary size, Encode strings and data protection)

Moderns Packers:

Hard to disassemble

polymorphic engines (uber-kewl)

VM and debugger's detection

time consuming

Millions are getting infected, thousands of new variants daily.

Polymorphic Code

Malware uses the same algorithm while its engine mutates the code.

The idea is to change the code each time the malware is being executed so it will create difficulties for AV vendors to produce efficient signatures as well as for malware researchers, OS Vendors (Most Are Lame) to understand the malware behavior.

Polymorphic Code

Start:

GOTO Decryption_Code Encrypted:

... lots of encrypted code ...

Decryption_Code:

A = Encrypted

Loop: B = *A

B = B XOR CryptoKey

*A = B

A = A + 1

GOTO Loop IF NOT A = Decryption_Code

GOTO Encrypted CryptoKey:

some_random_number

Polymorphic Code (Mutation)

Start:

GOTO Decryption_Code Encrypted: ->> same algorithm

... lots of encrypted code ... ->> same algorithm

Decryption_Code:

$c=2^{(c*4)} \ll \text{-- Mutation}$

A = Encrypted : ->> same algorithm

Loop: B = *A : ->> same algorithm

$C=C*(A+A) \ll \text{-- Mutation}$

B = B XOR CryptoKey

*A = B ->> same algorithm

$A = A + 1 \ll \text{-- Mutation}$

GOTO Loop IF NOT A = Decryption_Code

$C=C+8 \ll \text{-- Mutation}$

GOTO Encrypted CryptoKey:

some_random_number ->> same algorithm

Anti Debugging

some common ways malwares are using to check if a debugger is present.

.. Groovy Way ..

```
push offset @not_debugged
```

```
push dword fs:[0] <← Install SEH handler
```

```
mov fs:[0], esp <<- Set the Hook
```

```
push 1234h ;invalid handle <<- Generate Fault
```

```
call CloseHandle
```

```
; if fall here (cant MOV next), process is debugged
```

```
push @not_debugged
```

```
call SetUnhandledExceptionFilter <<- No Handler Should Exist
```

```
xor eax, eax
```

```
mov eax, dword [eax] ; trigger exception
```

```
;program terminated if debugged
```


Popular Packers

Armadillo – considered one of the best packers, supports encryption, debugger detection, vm detection.

Themida – considered one of the best packers, supports encryption, debugger detection

VMProtect – an advanced protection and compression product

UPX – Free and popular opensource packer, uses LZMA

ASPack – consider weak and outdated.

FSG – No loner maintained but still popular, supports only static PE

The Need Simplicity

- * Malware analysis is not an easy task.
- * Deep Knowledge with operating system internals.
- * Experience with low level programming.
- * Be a Bold Guy 😊

Why Automation ?

- * Not every one wants to disassemble code
- * Speed up the analysis process
- * Help malware researchers to test more variants and learn new attack vectors (Observing from the side)
- * its fun ?

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Coseinc Advanced Malware Analysis Lab



Goals

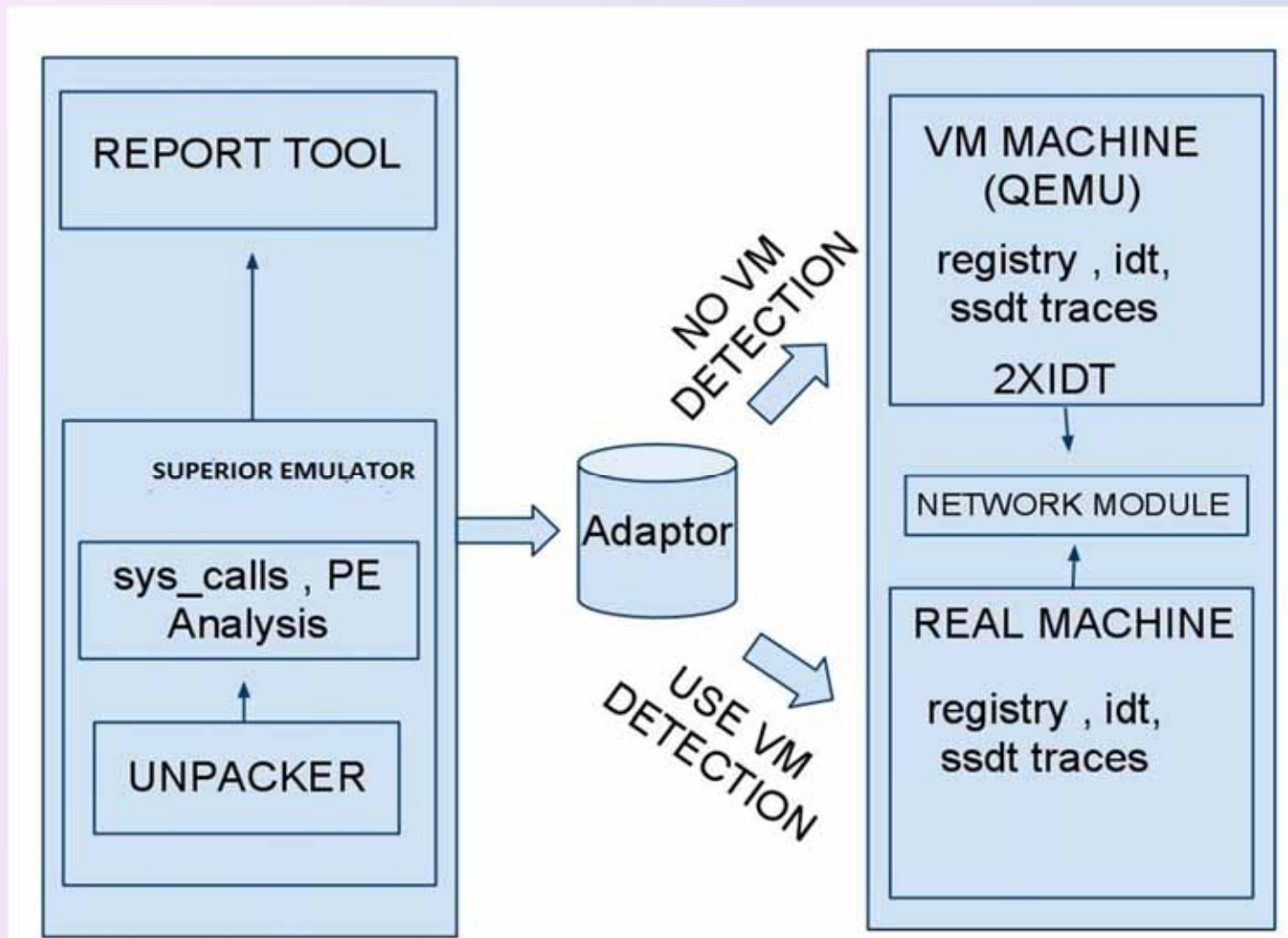
- *Research and learn new malwares attack vectors
- *Speed up the analysis process time
- *Provide the analyst with real time information regarding the binary activity inside the operating system internals



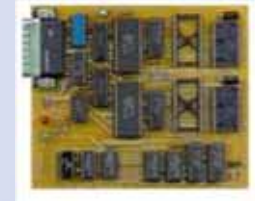
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CAMAL Internals

Framework Modules



Camal Superior Emulator (CSE)



*Emulating the following CPU architectures:

[ARM](#) | [MIPS](#) | [MOS 6502](#) | [Motorola 68K](#) | [Motorola 88K](#)
| [X86](#) | [PowerPC](#) | [FAPRA](#)

* build over low level virtual machine (LLVM) for the backend. user mode and system emulation, dynamic as well as static recompilation.

* Detect SIDT VMM query techniques by looking for ring 3 `0xffXXXXXX` address.

* 248K foot print.

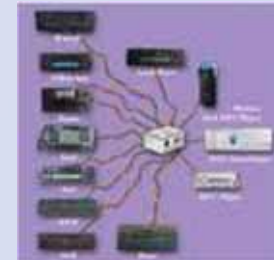
UNPACKER

- * Uses the Superior emulator for the execution process.
- * Support unpacking for the following packers:
 - UPX(x)
 - Aspack
 - PE-Compact
 - FSG
- * Support Real Time Brute Force (avoid packers lock).

PE Analysis

- * Section Header: The section table is an array of IMAGE_SECTION_HEADERS structures. An IMAGE_SECTION_HEADER provides information about its associated section, including location, length, and characteristics.
- * Import Tables : Load the imports the windows loader needs to use when executing, for instance kernel32.dll, user32.dll.
- * RVA: simply an offset in memory, relative to where the PE file was loaded. (PE GPS)

Adapter



- * Parse the emulator results
- * If vm detect found (SIDT 0xffff), send sample to real machine
- * Sample compiled for un-supported architecture
- * report hangs, signals, abnormal behavior

QEMU (Virtual machine)



- * pre-installed with CAMAL DTSR driver
- * compiled in emulation mode (non VMX support)
- * modified revert version



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Real Machine

- * Non emulation environment
- * Support reverting by reloading via BOOTP
- * Pre-installed with CAMAL DTSR driver

Software being used

- * QEMU *www.qemu.org*
- * tcpdump www.tcpdump.org
- * sysinternals *technet.microsoft.com/en-us/sysinternals/default.aspx*
- * debian *debian.org*



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LETS DEMO THE FRAMEWORK ...